1. Consider the case of three periodic tasks:

Task  $t_1$ :  $C_1 = 20 \text{ ms}; T_1 = 100 \text{ ms}$ Task  $t_2$ :  $C_2 = 40 \text{ ms}; T_2 = 150 \text{ ms}$ Task  $t_3$ :  $C_3 = 100 \text{ ms}; T_3 = 350 \text{ ms}$ 

Apply the Utilization Bound Theorem to determine if these tasks are schedulable using a rate monotonic scheduling strategy. Suppose the computation time for task 1 doubles to 40 msec, now determine if the tasks are schedulable, and then apply the less conservative Completion Time Theorem.

2. Suppose we have four tasks: two periodic, one aperiodic, and one interrupt driven aperiodic. The non-interrupt driven tasks require access to a shared data store, and we wish to give the interrupt-drive task the highest priority:

periodic task  $t_1$ :  $C_1 = 30$  ms,  $T_1 = 100$  ms aperiodic task  $t_2$ :  $C_2 = 30$  ms,  $T_2 = 150$  ms interrupt driven aperiodic task  $t_a$ :  $C_a = 10$  ms,  $T_a = 200$  ms periodic task  $t_3$ :  $C_3 = 30$  ms,  $T_3 = 300$  ms

The context switch time is included in the indicated CPU times. Use the Generalized Utilization Bound Theorem to determine if this task set is schedulable.

3. Given two tasks  $T_1$  and  $T_2$  with two shared data structures protected with binary semaphores  $S_1$  and  $S_2$ , show how the *priority ceiling protocol* prevents mutual deadlock and guarantees that a high-priority task will be blocked by at most one critical section of any lower priority task.